

FRACTAL

Certified Learning and Compilation of Finite State Automata

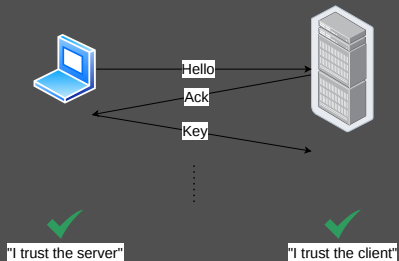
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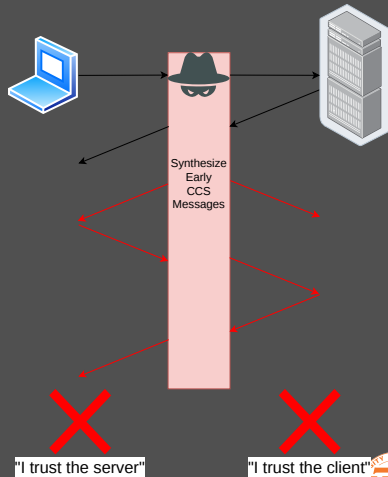
September 2026



TLS EarlyCCS Vulnerability

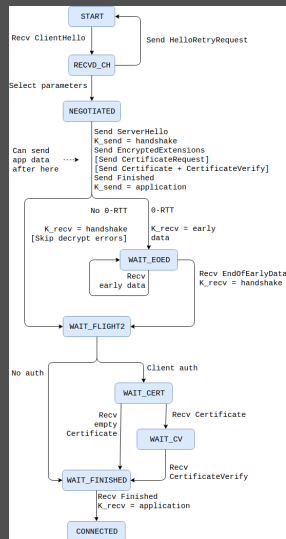


2-year vulnerability window for complete traffic decryption attacks.



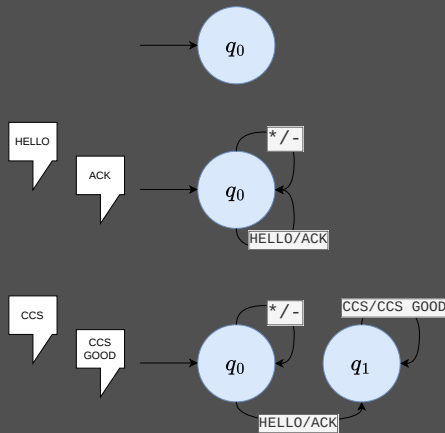
State Machines in Critical Systems

- **Finite State Machines** (FSMs) are a part of critical networking protocols, ICS, and hardware synthesis
- Very easy to implement and analyze state machines because of their limitations
- Because state machines are used in critical systems, we need to **verify** their **implementations** are accurate to **specifications**



Active Automata Learning

- **Learner** constructs state machines by asking the **Teacher** for **counterexamples**
- **Counterexample analysis** yields new states, transitions
- Iteratively improve system approximation until it is **relatively complete**
- Algorithms exist for DFAs, NFAs, Moore and Mealy machines, etc.



SOTA AAL Libraries

- **LearnLib** (Java, TU Dortmund)
 - L^* , KV, TTT, NL^* , and others for DFA, Mealy, and Moore machines
 - Used for state machine inference of TLS, SSH, TCP, and DTLS implementations, EMV bank cards, and industrial printer controllers
- **AALpy** (Python, TU Graz)
 - Active and passive learners for deterministic, non-deterministic, and stochastic automata
 - Used for Bluetooth LE and MQTT learning and device fingerprinting
- **LibALF** (C++, RWTH Aachen)
 - Active and passive learners (Angluin L^* , KV, NL^* , RPNI, Biermann)
 - Used for learning-based verification, e.g., inferring regular invariants



SOTA - Can you trust it?

■ LearnLib

- NL*: learned models contain spurious states (5.8 years)
- TTT: counterexample handling deviates from the algorithm, stalling refinement (1.9 years)

■ AALpy

- KV: counterexamples analyzed from the wrong system state, corrupting refinement (2.9 years)
- RPNl: incompatible states merged, producing wrong models (2.0 years)

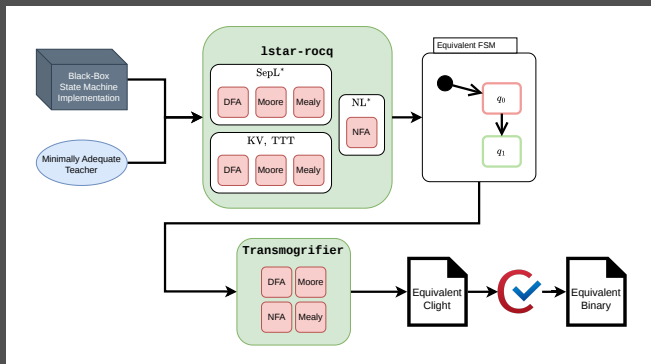
■ LibALF

- KV: hypotheses built with out-of-range states and misplaced accepting states (7 months)
- KV: rejecting states left without an output, yielding malformed hypotheses (8 months)
- L*: counterexamples silently dropped, stalling learning (7 months)



FRACTAL¹ - Certified AAL Implementations

- **First formally-verified** implementations of AAL algorithms
- Includes 4 algorithms supporting 4 FSM classes
- **Guaranteed** to return **correct**, **minimal** FSMs given a Teacher
- Extracts to performant **OCaml**



¹Framework in Rocq for **ACTIVE** Automata Learning

Challenges in Certified AAL

- AAL algorithms are usually given in an **imperative** style, which is usually much more difficult to represent and verify than a **functional**-style program
- Imperative programs require **termination arguments**, which are occasionally underspecified in AAL papers
- Extracted OCaml code is often **not performant** by default

So the goal is: implement and verify **imperative-style** algorithms consisting of **complex termination arguments** that we can extract to **performant native code**. Sounds like a lot of work!



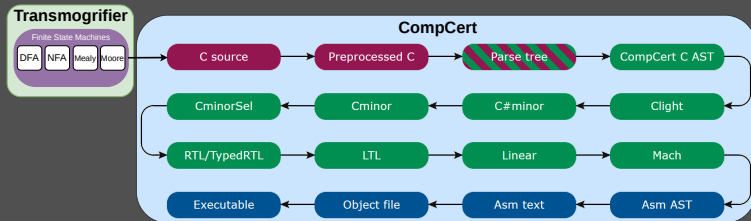
A lot of work - Learning

- Built in the **Rocq** theorem prover environment
- Uses a common **fuel pattern** to represent imperative programs as functional, cleanly folding into explicit **termination arguments**
- Supports L^* , TTT, KV, and NL^* for DFAs, NFAs, Moore machines, and Mealy machines
- Implementations are optimized to approach the performance of equivalent C code
- Can attain even better performance by combining with...



A lot more work - Compilation

- Extracted OCaml FSM code is pretty fast... but machine code compiled from C is usually a lot faster
- **Transmogrifier** is our **formally-verified FSM compiler**
- Converts FSMs into **Clight** code, a subset of C for the **CompCert** certified C compiler
- Guaranteed to generate functionally-equivalent machine code



Transmogrier Results

```

const state_t table[|Q|·|Σ|] = {(δ(qi, Σj))(|Q|,|Σ|)};
const output_t atable[|Q|] = {(qi ∈ F)|Q|};
const state_t q0 = q0;

state_t delta(state_t q, symbol_t a) { // δ(q,a)
  if (q < |Q| & a < |Σ|)
    return table[q * |Σ| + a];
  else
    return |Q|; // Invalid state index
}

output_t accept(state_t q) { // q ∈ F
  if (q < |Q|)
    return atable[q];
  else return 2; // Neither true nor false
}

state_t run(word_t w, size_t len) { // δ*(q0, w)
  size_t i = 0; state_t q = q0;
  while (i < len) {
    q = delta(q, w[i++]);
  }
  return q;
}

```

Figure 4. Shape of generated Clight code for DFAs

Table 3. Comparison of FSMs in OCaml (OC) and CompCert-compiled C on strings of length 10^6 averaged over 100 runs. Memory usage is the maximum Resident Set Size throughout all runs. Improvement (Imp.) is given by $\frac{\text{OCaml}}{\text{C}}$.

Test	Runtime (ms)			Mem. Usage (MB)		
	OC	C	Imp.	OC	C	Imp.
dfa.alternating	42.9	2.8	15.2×	30.7	9.5	3.23×
dfa.mod3	42.7	2.8	15.3×	30.8	9.5	3.24×
dfa.div7	134.7	3.1	44.0×	31.1	9.3	3.34×
nfa.suffix	191.0	2.8	68.4×	43.8	9.3	4.71×
moore.traffic	41.2	2.5	16.4×	30.6	9.4	3.26×
mealy.vending	301.6	3.0	100.6×	109.4	17.2	6.36×

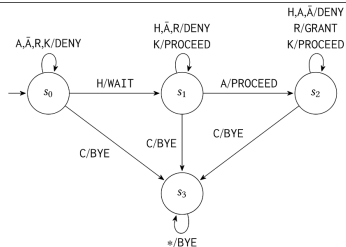


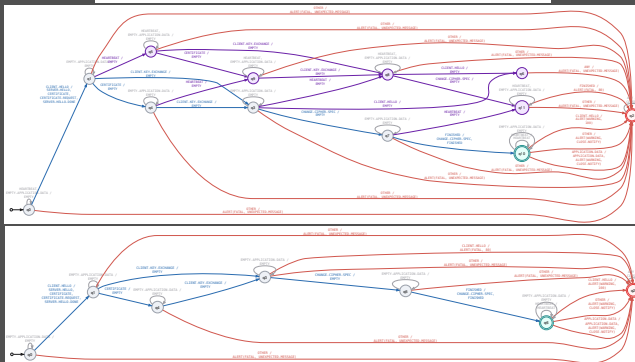
Figure 5. The learned session handshake machine. s_0 is initial, s_1 is reached once HELLO is seen, s_2 once AUTH succeeds, and s_3 absorbs after CLOSE. Edges are labelled *input/output*.



FRACTAL Results

Table 2. Membership query comparisons between FRACTAL and LearnLib (LL) on the LearnLib integration test suite with query caching

Algorithm	FRACTAL	LL	Overhead
L*	6760	6018	12%
KV	2681	2335	15%
TTT	3805	3494	9%
NL*	268	258	4%



Future work

- **More support - Buchi automata** via coinduction, register automata, more algorithm variants, etc.
- **Optimization** - reduce the query complexity gap between FRAC TAL and LearnLib
- **Discovery** - apply FRAC TAL to more real-world systems to discover state machine vulnerabilities

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